

# An Efficient Design of Reconfigurable Reversible Gate for Data Encryption and Decryption

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**Abstract**— The development in the field of nanometer technology leads to minimize the power consumption of logic circuits. Reversible logic design has been one of the promising technologies gaining greater interest due to less dissipation of heat and low power consumption. On the other hand, reversible logic circuits can decrease energy dissipation theoretically to zero. Recently an approach to encryption/decryption based on using reversible logic circuits has been proposed. This paper presents a solution to designing encryption and decryption schemes based entirely on reversible logic. In our solution a building block of an encryption and decryption scheme is a cascade of 4-input reversible gates. In this way the building block can perform any reversible 4-variable function. For this purpose a reconfigurable reversible gate has been proposed. The design of such a reconfigurable gate built from standard reversible gates, i.e. NOT, CNOT, Toffoli and Fredkin gates, is presented. In the paper a complete scheme for encryption/decryption of 9-bit data is described using Verilog HDL language and its quantum cost is calculated. Simulation and verification of this scheme is presented on Xilinx ISE 14.7.

**Key words:** Encryption, Decryption, Reversible Logic Circuits, Reconfigurable Reversible Gate

## I. INTRODUCTION

The growing technologies have increased the demand of high performance computing. According to G. Moore's law [1], number of transistor counts to be integrated per unit area in devices will almost double in one and half year. To achieve high speed computation, high packaging density in the logic circuits is required which results in more heat dissipation. The conventional computing is found unable to deal with low power, high compaction and heat dissipation issues of the current computing environment. Recently, it has been applied to cryptography [2-10]. A circuit (gate) is called reversible if there is a one-to-one correspondence between its inputs and outputs. Research on reversible logic circuits is motivated by advances in quantum computing, nanotechnology and low power design. Therefore, reversible logic synthesis has been intensively studied recently. The attention is focused mainly on the synthesis of circuits built from the NCT library of gates, i.e. NOT, CNOT and Toffoli gates. Modern simulation tools based on FPGAs have enabled modeling of such circuits [11]. The design of a cipher in VHDL was described [9]. In the paper we study an application of reversible logic to developing encryption circuits. The simple implementation of a cipher using reversible circuits was the aim of this work. Each gate used in a cascade of reversible gates is determined by the main key. By choosing different main keys different cascades and different substitution encryption are determined. For this purpose a reconfigurable reversible gate has been proposed. The design of this gate is presented for the first time in the literature. Results of Xilinx ISE based

simulation of a simple encryption circuit built from reversible gates are also presented.

The paper is organized as follows. Section 2 recalls basic concepts of reversible logic. In Section 3 design of reconfigurable reversible gate (RRG) is introduced. In Section 4 an implementation of encryption/decryption by cascades of RRG gates is given. In Section 5 simulation of the cipher in Xilinx ISE is presented. Section 6 summarizes the paper.

## II. REVERSIBLE DEFINITIONS

Let us recall basic notions of reversible logic gates and circuits.

### A. Definition 1

A completely specified n-input n-output Boolean function (referred to as n\*n function) is called reversible if it maps each input assignment into a unique output assignment. There are  $2^n!$  Reversible n\*n Boolean functions. For n = 3 this number equals to 40,320, for n = 4 is greater than  $2^{1013}$ .

### B. Definition 2

An n-input n-output (n\*n) gate (or circuit) is reversible if it realizes an n\*n reversible function. In a reversible circuit fan out of each gate output is always equal to 1. As a consequence n\*n reversible circuits can be only built as a cascade of k\*k reversible gates ( $k \leq n$ ).

### C. Definition 3

A set of reversible gates that can be used to build reversible circuits is called a gate library.

Many gate libraries have been examined in the literature. The so called NCT library for  $n \leq 4$  consists of 1\*1 NOT, 2\*2 CNOT and 3\*3 and 4\*4 TOFFOLI gates. Definitions of mixed polarity generalizations of these gates as well as 3\*3 Fredkin gate are given below.

### D. Definition 4

Let  $a_i \in \{0, 1\}$  for  $i = 1, 2, 3, 4$ , and let  $x'$  denotes negation of  $x$ .

1\*1 NOT( $x_1$ ) gate performs the operation  $(x_1) - (x_1 \oplus 1)$ ,  
2\*2 CNOT( $x_1, x_2$ ) gate performs the operation  $(x_1, x_2) - (x_1, (x_1 \oplus a_1) \oplus x_2)$ ,

3\*3 TOFFOLI( $x_1, x_2, x_3$ ) gate performs the operation  $(x_1, x_2, x_3) - (x_1, x_2, (x_1 \oplus a_1)(x_2 \oplus a_2) \oplus x_3)$ ,

4\*4 TOFFOLI4( $x_1, x_2, x_3, x_4$ ) gate performs the operation  $(x_1, x_2, x_3, x_4) - (x_1, x_2, x_3, (x_1 \oplus a_1)(x_2 \oplus a_2)(x_3 \oplus a_3) \oplus x_4)$ ,

3\*3 FREDKIN( $x_1, x_2, x_3$ ) gate performs the operation  $(x_1, x_2, x_3) - (x_1, x_1' x_2 + x_1 x_3, x_1 x_2 + x_1' x_3)$ .

The above defined  $i^*i$  gates, where  $i = 1, 2, 3, 4$ , (in short, denoted by N, C, T, T4 and F, respectively) invert input  $x_i$  if and only if the values of inputs  $x_1, x_2, \dots, x_{i-1}$  differ from corresponding  $a_1, a_2, \dots, a_{i-1}$  coefficients, passing these inputs unchanged to corresponding outputs. Signals which

are passed unchanged from input to output of the gate are called control lines. The signal  $x_i$  which can be modified by the gate is called target. In the Fredkin gate  $x_1$  is the control line. When its input signal is 0 then the signals  $x_2$  and  $x_3$  are passed to outputs unchanged, and when it's input signal is 1 then the signals  $x_2$  and  $x_3$  are swapped.

Some commonly used names for describing control lines have been introduced. If all  $a_1, a_2, \dots, a_i$  coefficients equals zero, the gate has positive-polarity control lines. Alternatively, if all  $a_1, a_2, \dots, a_i$  coefficients equals one, the gate has negative-polarity control lines. These names come from expanding the expressions for the functions realized by the gate. Positive-polarity means that all inputs which correspond to control lines directly affect the target line. Negative-polarity means that the target line is affected only if the values of control lines are equal to 0. The term mixed-polarity control lines is used if all values of the coefficients  $a_1, a_2, \dots, a_i$  are allowed to be either 0 or 1. Pictorial symbols of the above defined positive-polarity gates are shown in Fig. 1, and the pictorial symbol of Fredkin gate is shown in Fig. 2.

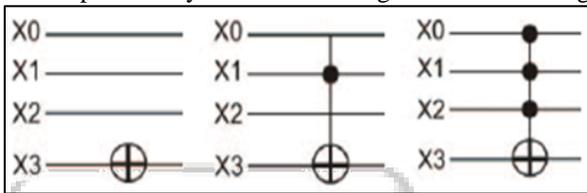


Fig. 1: Graphical representations of reversible gates N3, C3-1 and T3-210.

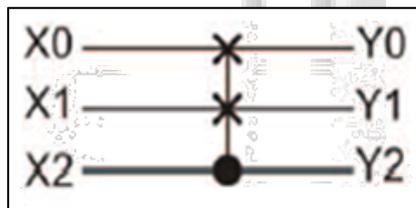


Fig. 2: Graphical representation of Fredkin gate F10-2.

Pictorial symbols of mixed-polarity gates are shown in Fig. 4 (gates B3, B4, B5, B7, B8 and B9). In the description of a gate from NCT library, e.g. T3-210 (see Fig. 1), the first number appearing after the gate's name (in this case: T) denotes the target line (3) while the numbers appearing after the hyphen denote control lines (2, 1, 0). In the description of a Fredkin gate, e.g. F10-2, the first two numbers appearing after the gate's name (in this example: F) denote controlled (swapping) lines (1 and 0) while the number appearing after the hyphen denotes control line (2). The number of different positive-polarity gates belonging to NCT library which can appear in 4-line circuits is 32 (4 NOT gates, 12 CNOT gates, 12 Toffoli gates and 4 T4 gates).

### III. IMPLEMENTATION OF RRG

Reconfigurable reversible gate (RRG) replaces one of 32 gates from the NCT library in the cipher. This is why besides 4 inputs transferring information data there have to be five lines for choosing one of 32 types of gates. At the present state-of-the-art of reversible logic synthesis it is not possible to optimally synthesize 9-input reversible gate. After considering many variants of gates RRG the circuit shown in Fig. 3 has been constructed with the relatively small quantum cost equal to 79.

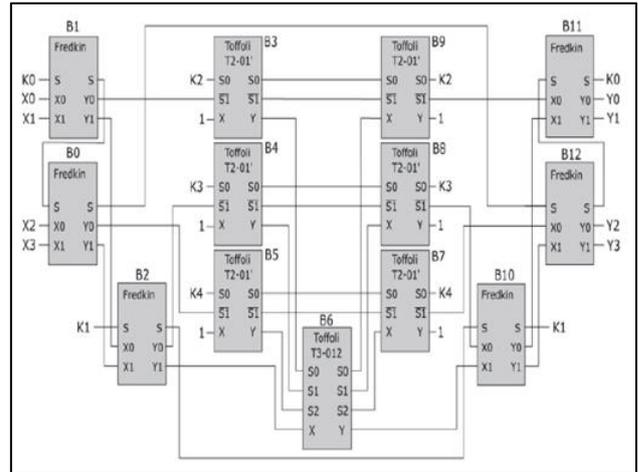


Fig. 3: General scheme of Reconfigurable Reversible Gate (RRG).

The first three Fredkin gates B0, B1 and B2 select the input signal which will be modified by the 4- input Toffoli gate B6. The next three 3-input Toffoli gates B3, B4 and B5 feed control signals for the gate B6 that are selected by the given configuration either input or constant signals. Thus functioning of the gate B6 is determined by signals K. Gates B7, B8 and B9 reconstruct initial constant signals, while gates B10, B11 and B12 fix the order of output signals  $Y_3, Y_2, Y_1$  and  $Y_0$ . The RRG built from basic reversible gates is shown in Fig. 4.

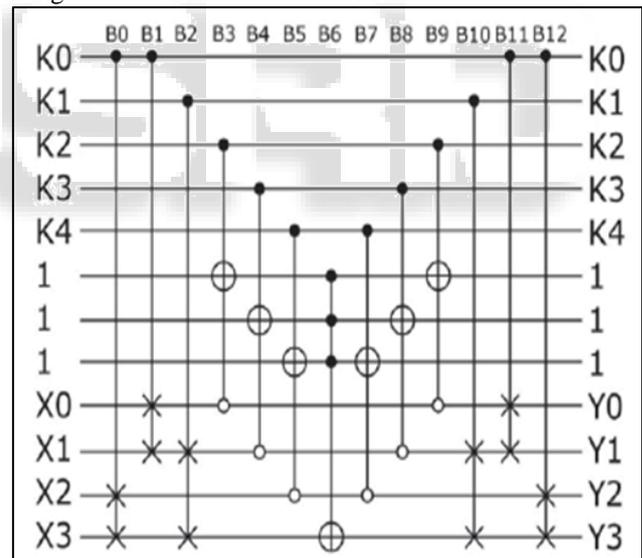


Fig. 4: Reconfigurable Reversible Gate (RRG)

For simulation of RRG a procedure B\_FREDKIN for representing functioning of Fredkin gate and a procedure B\_NCT for representing any 4-input gate from NCT library have been constructed in VHDL and shown in Tables I and II.

```

procedure B_FREDKIN
(X : in std_logic_vector; -- input line
 C : in integer; -- number of control line
 L1 : in integer; -- number of first line
 L2 : in integer; -- number of second line
 Y : out std_logic_vector) -- output line
is
begin
if X(C)='0' then Y:=X;
else Y:=X; Y(L1):=X(L2);
Y(L2):=X(L1);
end if;
end procedure B_FREDKIN;

```

Table 1: Procedure B\_Fredkin

In the procedure B\_NCT signals C1, C2 and C3 determine number of line to which control inputs of Toffoli gate are attached while the number of controlled line is determined by signal C0. Binary signals N1, N2 and N3 determine polarization of the control signal Ci. Value N1 = 1 indicates that i-th signal is a negative control signal.

```

procedure B_NCT
(X : in std_logic_vector; -- input line
 C1 : in integer; -- control line
 N1 : in std_logic; -- negative control line C1
 C2 : in integer; -- control line
 N2 : in std_logic; -- negative control line C2
 C3 : in integer; -- control line
 N3 : in std_logic; -- negative control line C3
 C0 : in integer; -- control line output
 Y : out std_logic_vector) -- output line
begin
if C1=12 then Y:=X; Y(C0) := not X(C0);
elsif C2=12 then Y:=X;
Y(C0) := (X(C1) xor N1) xor X(C0);
elsif C3=12 then Y:=X;
Y(C0) := ((X(C1) xor N1)
and (X(C2) xor N2))xor X(C0);
else Y:=X;
Y(C0) := ((X(C1) xor N1)
and (X(C2) xor N2)
and (X(C3) xor N3))xor X(C0);
end if;
end procedure B_NCT;

```

Table 2: Procedure B Nct

The above presented procedures were used for describing the gate RRG as shown below.  
 $KX(0) := (KEY(0) \& KEY(1) \& KEY(2) \& KEY(3) \& KEY(4) \& "111" \& DIN(0) \& DIN(1) \& DIN(2) \& DIN(3));$   
B\_FREDKIN (KX( 0),0,10,11, KX( 1)); B\_FREDKIN (KX( 1),0, 8, 9, KX( 2));  
B\_FREDKIN (KX( 2),1, 9,11, KX( 3)); B\_NCT (KX(3),2,'0', 8,'1',12,'0', 5,KX( 4));  
B\_NCT (KX(4),3,'0', 9,'1',12,'0', 6,KX( 5)); B\_NCT (KX(5),4,'0',10,'1',12,'0', 7,KX( 6));  
B\_NCT (KX(6),5,'0', 6,'0', 7,'0',11,KX( 7)); B\_NCT (KX(7),4,'0',10,'1',12,'0', 7,KX( 8));  
B\_NCT (KX(8),3,'0', 9,'1',12,'0', 6,KX( 9)); B\_NCT (KX(9),2,'0', 8,'1',12,'0', 5,KX(10));  
B\_FREDKIN (KX(10),1, 9,11, KX(11)); B\_FREDKIN (KX(11),0, 8, 9, KX(12));  
B\_FREDKIN (KX(12),0,10,11, KX(13));  
DOUT(3 downto 0)<=(KX(13)(11) & KX(13)(10) & KX(13)(9) & KX(13)(8));  
12-bit word KX is the state of line just before or just after gate and is transferred from i-th to (i+1)-th gate.

#### IV. DESIGN OF ENCRYPTION BY RRG GATE

General idea of a cipher built from reversible gates was presented in [9]. It is shown in Fig. 5. The basic element of the cipher is a cascade of 16 4-input reversible gates. The same main key is used for encryption and decryption. The order of gates in the cascade for decryption is reversed in comparison with the cascade for encryption what ensures that it transforms ciphertext into plaintext.

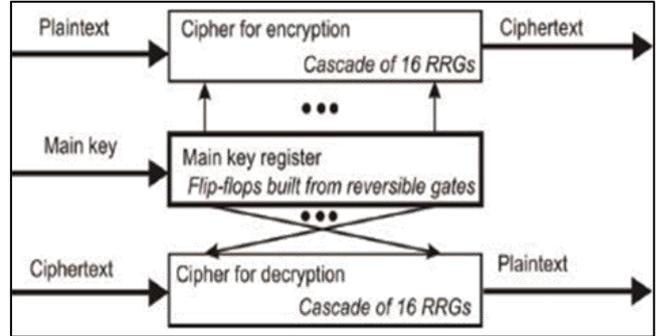


Fig. 5: General idea of a cipher for encryption and decryption built from reversible gates

It has been proved in [12-13] that optimal circuits for any of the 16! (equal to over  $2 \times 10^{13}$ ) require at most 15 4-input gates. Thus the circuit shown in Fig. 6 enables realization of any 4-variable reversible function. The circuit has 80 inputs denoted K (with indexes) that are partitioned into groups with 5 inputs in each of them. A 5-line group  $K[(5 \cdot (i+1) - 1) : 5 \cdot i]$  is used to configure i-th RRG gate. All inputs K are transferred to outputs so they can be reused for controlling the next gate. There are 4 data inputs X[3:0] and three lines with constant inputs (equal to 1) and identical outputs. Detailed description of the main key register as well as the circuit modifying its contents during encryption/decryption are presented in [14].

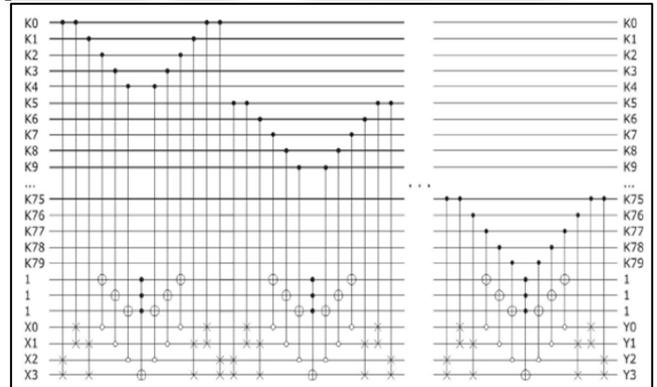


Fig. 6: The cascade of reversible gates implementing the 4-bit cipher.

#### V. SYNTHESIS AND SIMULATION RESULTS

For encryption and decryption of 5-bit key and two 4-bit ciphers and the main key register were used. In this section, first we will see the synthesis and simulation of the Encryption and Decryption using reconfigurable reversible gate. Encryption and Decryption using reconfigurable reversible gates are designed on Xilinx ISE 14.7 with Verilog HDL. The RTL schematics and simulation results of the Proposed Reversible Encryption and Decryptions are shown below.

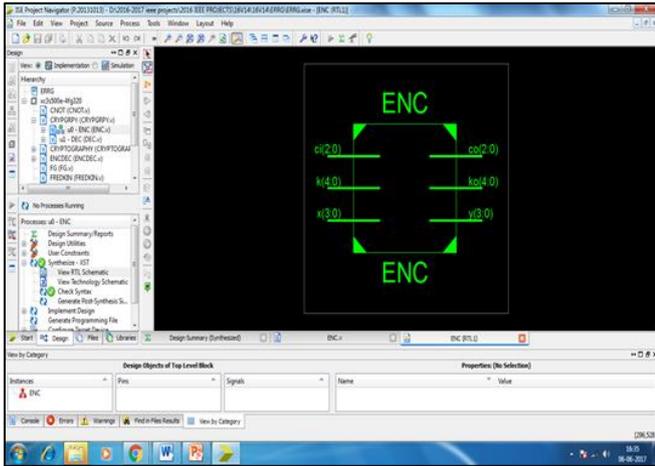


Fig. 7: RTL Schematic of Proposed Reversible Encryption

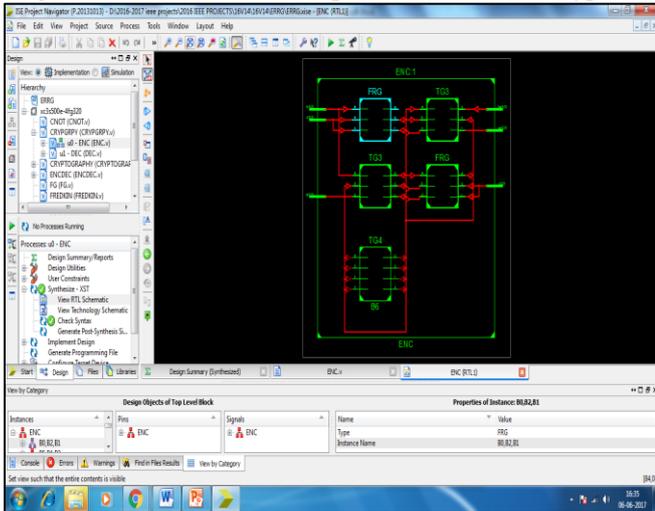


Fig. 8: Internal RTL Schematic of Proposed Reversible Encryption

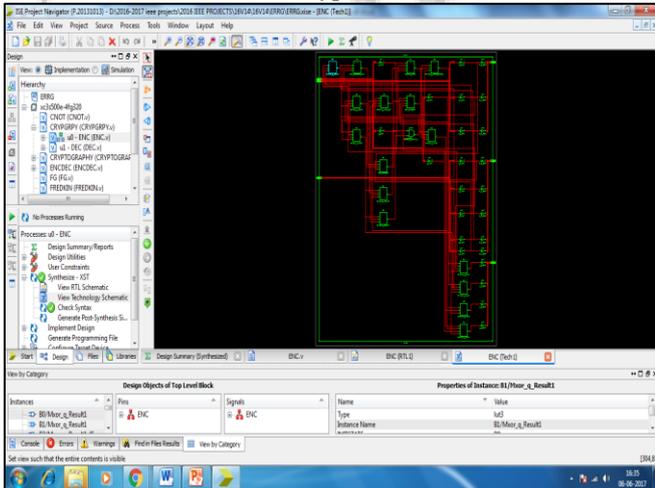


Fig. 9: Technology Schematic of Proposed Reversible Encryption

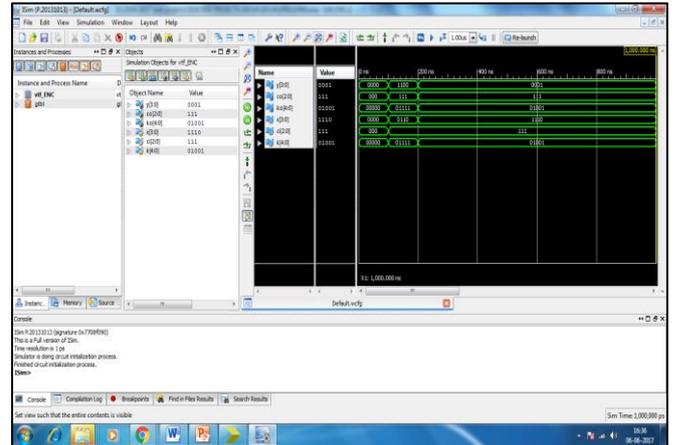


Fig. 10: Simulation of Proposed Reversible Encryption

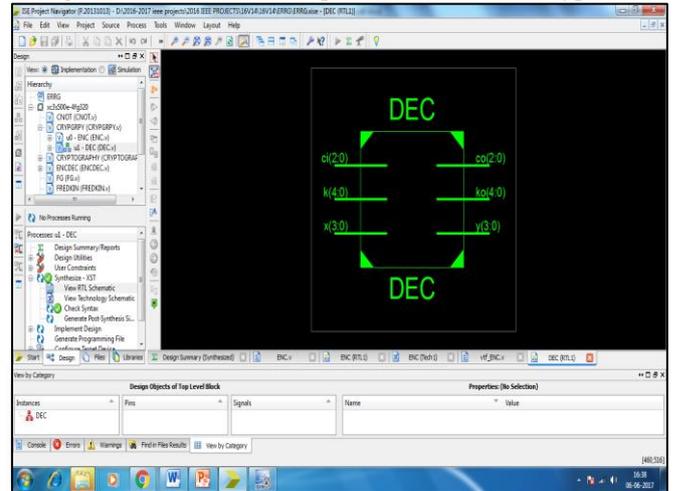


Fig. 11: RTL Schematic of Proposed Reversible Decryption

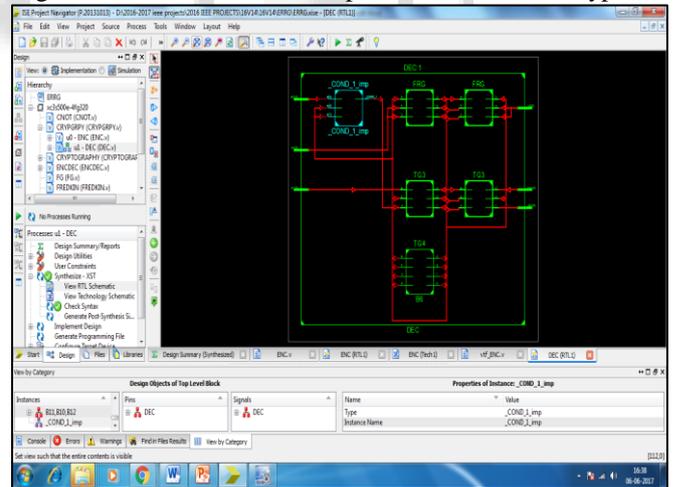


Fig. 12: Internal RTL Schematic of Proposed Reversible Decryption

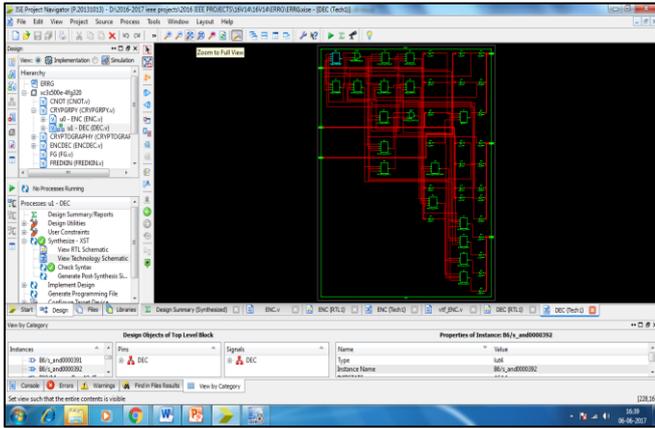


Fig. 13: Technology Schematic of Proposed Reversible Decryption

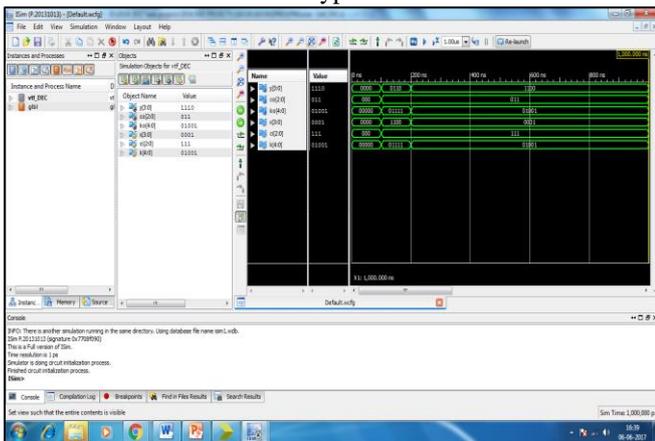


Fig. 14: Simulation of Proposed Reversible Decryption

DESIGN	GC	CI	GO	QC
ENCRYPTION	13	3	0	79
DECRYPTION	13	3	0	79

Table 3: Rlg Factors Comparison for Encryption and Decryption

Note: Gate count (GC), Garbage Outputs (GO), Constant Input (CI) and Quantum Cost (QC)

## VI. CONCLUSIONS

Firstly we have coded Verilog code for the reversible data encryption and decryption and all the synthesis and simulation results are implemented on Xilinx ISE 14.7. The main aim of the paper is a design of simple reconfigurable reversible gate (RRG) which enables implementation of any of the 32 4-input reversible gates from the NCT library. The design of RRG is presented for the first time in the literature. An application of this gate to implement ciphers for encryption and decryption of binary data is described. Results of encryption and decryption simulation of the cipher built from reversible gates are also presented.

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